## The Solution of the Kadison-Singer Problem

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## Outline

Disclaimer

The Kadison-Singer Problem, defined.
Restricted Invertibility, a simple proof.
Break
Kadison-Singer, outline of proof.

## The Kadison-Singer Problem (‘59)

A positive solution is equivalent to:
Anderson's Paving Conjectures ('79, ‘81)
Bourgain-Tzafriri Conjecture ('91)
Feichtinger Conjecture ('05)
Many others

Implied by:
Akemann and Anderson’s Paving Conjecture ('91)
Weaver's KS 2 Conjecture

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## The Kadison-Singer Problem ('59)

Let $\mathcal{A}$ be a maximal Abelian subalgebra of $\mathcal{B}\left(\ell^{2}(\mathbb{N})\right)$, the algebra of bounded linear operators on $\ell^{2}(\mathbb{N})$

Let $\rho: \mathcal{A} \rightarrow \mathbb{C}$ be a pure state.
Is the extension of $\rho$ to $\mathcal{B}\left(\ell^{2}(\mathbb{N})\right)$ unique?

See Nick Harvey's Survey or Terry Tao’s Blog

## Anderson's Paving Conjecture ‘79

For all $\epsilon>0$ there is a $k$ so that for every $n$-by-n symmetric matrix $A$ with zero diagonals,
there is a partition of $\{1, \ldots, n\}$ into $S_{1}, \ldots, S_{k}$

$$
\left\|A\left(S_{j}, S_{j}\right)\right\| \leq \epsilon\|A\| \quad \text { for } \quad j=1, \ldots, k
$$

Recall $\|A\|=\max _{\|x\|=1}\|A x\|$

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Recall $\|A\|=\max _{\|x\|=1}\|A x\|$


## Anderson's Paving Conjecture ‘79

For all $\epsilon>0$ there is a $k$ so that for every self-adjoint bounded linear operator $A$ on $\ell_{2}$,
there is a partition of $\mathbb{N}$ into $S_{1}, \ldots, S_{k}$

$$
\begin{aligned}
& \left\|A\left(S_{j}, S_{j}\right)\right\| \leq \epsilon\|A\| \quad \text { for } \quad j=1, \ldots, k \\
& \|A\|=\sup _{\|x\|=1}\|A x\|
\end{aligned}
$$

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$$
\left\|A\left(S_{j}, S_{j}\right)\right\| \leq \epsilon\|A\| \quad \text { for } \quad j=1, \ldots, k
$$

Is equivalent if restrict to projection matrices.
[Casazza, Edidin, Kalra, Paulsen ‘07]

## Anderson's Paving Conjecture ‘79

Equivalent to [Harvey '13]:

There exist an $\epsilon>0$ and a $k$ so that for $v_{1}, \ldots, v_{n} \in \mathbb{C}^{d}$
such that $\left\|v_{i}\right\|^{2} \leq 1 / 2$ and $\sum v_{i} v_{i}^{*}=I$
then exists a partition of $\{1, \ldots, n\}$ into $k$ parts s.t.

$$
\left\|\sum_{i \in S_{j}} v_{i} v_{i}^{*}\right\| \leq 1-\epsilon
$$

## Moments of Vectors

The moment of vectors $v_{1}, \ldots, v_{n}$ in the direction of a unit vector $u$ is $\sum_{i}\left(v_{i}^{T} u\right)^{2}$
$v_{1}$

1

2.5
$v_{1}$

4

## Moments of Vectors

The moment of vectors $v_{1}, \ldots, v_{n}$ in the direction of a unit vector $u$ is

$$
\begin{aligned}
\sum_{i}\left(v_{i}^{T} u\right)^{2} & =\sum_{i} u^{T}\left(v_{i} v_{i}^{T}\right) u \\
& =u^{T}\left(\sum_{i} v_{i} v_{i}^{T}\right) u
\end{aligned}
$$

## Vectors with Spherical Moments



For every unit vector $u$

$$
\sum_{i}\left(v_{i}^{T} u\right)^{2}=1
$$

## Vectors with Spherical Moments



For every unit vector $u$

$$
\begin{aligned}
& \sum_{i}\left(v_{i}^{T} u\right)^{2}=1 \\
& \sum_{i} v_{i} v_{i}^{T}=I
\end{aligned}
$$

Also called isotropic position

# Partition into Approximately $1 / 2$-Spherical Sets 




## Partition into Approximately $1 / 2$-Spherical Sets



$$
1 / 4 \leq \sum_{i \in S_{j}}\left(v_{i}^{T} u\right)^{2} \leq 3 / 4
$$

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$$

$$
1 / 4 \leq \operatorname{eigs}\left(\sum_{i \in S_{j}} v_{i} v_{i}^{T}\right) \leq 3 / 4
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$$

$$
\Longleftrightarrow \operatorname{eigs}\left(\sum_{i \in S_{j}} v_{i} v_{i}^{T}\right) \leq 3 / 4
$$

because $\sum_{i \in S_{1}} v_{i} v_{i}^{T}=I-\sum_{i \in S_{2}} v_{i} v_{i}^{T}$

## Partition into Approximately $1 / 2$-Spherical Sets





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$$
\Longleftrightarrow \quad\left\|\sum_{i \in S_{j}} v_{i} v_{i}^{T}\right\| \leq 3 / 4
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## Big vectors make this difficult



## Big vectors make this difficult



## Weaver's Conjecture $\mathrm{KS}_{2}$

There exist positive constants $\alpha$ and $\epsilon$ so that
if all $\left\|v_{i}\right\| \leq \alpha$
then exists a partition into $S_{1}$ and $S_{2}$ with

$$
\operatorname{eigs}\left(\sum_{i \in S_{j}} v_{i} v_{i}^{*}\right) \leq 1-\epsilon
$$

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Implies Akemann-Anderson Paving Conjecture, which implies Kadison-Singer

## Main Theorem

For all $\alpha>0$
if all $\left\|v_{i}\right\| \leq \alpha$
then exists a partition into $S_{1}$ and $S_{2}$ with

$$
\operatorname{eigs}\left(\sum_{i \in S_{j}} v_{i} v_{i}^{*}\right) \leq \frac{1}{2}+3 \alpha
$$

Implies Akemann-Anderson Paving Conjecture, which implies Kadison-Singer

## A Random Partition?

Works with high probability if all $\left\|v_{i}\right\|^{2} \leq O(1 / \log d)$ (by Tropp '11, variant of Matrix Chernoff, Rudelson)

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\text { are } 1 / \alpha^{2} \text { of each }
$$



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Chance there exists a direction in which all land in same set is

$$
1-\left(1-2^{-1 / \alpha^{2}}\right)^{d} \rightarrow 1
$$

## The Graphical Case

From a graph $G=(V, E)$ with $|V|=n$ and $|E|=m$ Create $m$ vectors in $n$ dimensions:

$$
v_{a, b}(c)= \begin{cases}1 & \text { if } c=a \\ -1 & \text { if } c=b \\ 0 & \text { otherwise }\end{cases}
$$

$$
\sum_{(a, b) \in E} v_{a, b} v_{a, b}^{T}=L_{G}
$$

If G is a good d -regular expander, all eigs close to d very close to spherical

## Partitioning Expanders

Can partition the edges of a good expander to obtain two expanders.

Broder-Frieze-Upfal '94:
construct random partition guaranteeing degree at least d/4, some expansion

Frieze-Molloy '99: Lovász Local Lemma, good expander

Probability is works is low, but can prove non-zero

## Interlacing Families of Polynomials

A new technique for proving existence from very low probabilities

## Restricted Invertibility (Bourgain-Tzafriri)

## Special case:

For $v_{1}, \ldots, v_{n} \in \mathbb{C}^{d}$ with $\sum_{i} v_{i} v_{i}^{*}=I$
for every $k \leq d$ there is a $S \subset\{1, \ldots, n\},|S|=k$
so that

$$
\lambda_{k}\left(\sum_{i \in S} v_{i} v_{i}^{*}\right) \geq\left(1-\sqrt{\frac{k}{d}}\right)^{2} \frac{d}{n}
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Is far from singular on the span of $\left\{v_{i}\right\}_{i \in S}$

## Restricted Invertibility (Bourgain-Tzafriri)

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so that $\quad \lambda_{k}\left(\sum_{i \in S} v_{i} v_{i}^{*}\right) \geq\left(1-\sqrt{\frac{k}{d}}\right)^{2} \frac{d}{n}$

$$
\begin{aligned}
& \text { For } k=1 \text { says } \lambda_{1}\left(v v^{*}\right) \gtrsim \frac{d}{n} \\
& \text { while } \lambda_{1}\left(v v^{*}\right)=v^{*} v=\|v\|^{2} \approx \frac{d}{n}
\end{aligned}
$$

Similar bound for $k$ a constant fraction of $d!$

## Method of proof

Let $r_{1}, \ldots, r_{k}$ be chosen uniformly from $\left\{v_{1}, \ldots, v_{n}\right\}$

1. $\mathbb{E} \chi\left[\sum r_{j} r_{j}^{*}\right](x)$ is real rooted
the characteristic polynomial in the variable $x$ of the matrix inside the brackets

## Method of proof

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1. $\mathbb{E} \chi\left[\sum r_{j} r_{j}^{*}\right](x)$ is real rooted
2. $\lambda_{k}\left(\mathbb{E} \chi\left[\sum r_{j} r_{j}^{*}\right](x)\right) \geq\left(1-\sqrt{\frac{k}{d}}\right)^{2} \frac{d}{n}$
the $k$-th root of the polynomial

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2. $\lambda_{k}\left(\mathbb{E} \chi\left[\sum r_{j} r_{j}^{*}\right](x)\right) \geq\left(1-\sqrt{\frac{k}{d}}\right)^{2} \frac{d}{n}$
3. With non-zero probability

$$
\lambda_{k}\left(\chi\left[\sum r_{j} r_{j}^{*}\right](x)\right) \geq \lambda_{k}\left(\mathbb{E} \chi\left[\sum r_{j} r_{j}^{*}\right](x)\right)
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Because is an interlacing family of polynomials

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## Rank-1 updates of characteristic polynomials

$$
\text { As } \sum_{i} v_{i} v_{i}^{*}=I, \quad \mathbb{E} r_{j} r_{j}^{*}=\frac{1}{n} I
$$

Lemma: For a symmetric matrix $A$,

$$
\mathbb{E} \chi\left[A+r_{j} r_{j}^{*}\right](x)=\left(1-\frac{1}{n} \partial_{x}\right) \chi[A](x)
$$

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$$

Proof: follows from rank-1 update for determinants:

$$
\operatorname{det}\left(A+u u^{*}\right)=\operatorname{det}(A)\left(1+u^{*} A^{-1} u\right)
$$

## The expected characteristic polynomial

Lemma: For a symmetric matrix $A$,

$$
\mathbb{E} \chi\left[A+r_{j} r_{j}^{*}\right](x)=\left(1-\frac{1}{n} \partial_{x}\right) \chi[A](x)
$$

Corollary:

$$
\mathbb{E} \chi\left[\sum_{j=1}^{k} r_{j} r_{j}^{*}\right](x)=\left(1-\frac{1}{n} \partial_{x}\right)^{k} x^{d}
$$

## Real Roots

Lemma: if $p(x)$ is real rooted, so is $\left(1-c \partial_{x}\right) p(x)$


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$+$

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$$
\mathbb{E} \chi\left[\sum_{j=1}^{k} r_{j} r_{j}^{*}\right](x)=\left(1-\frac{1}{n} \partial_{x}\right)^{k} x^{d}
$$

So, $\mathbb{E} \chi\left[\sum_{j=1}^{k} r_{j} r_{j}^{*}\right](x)$ is real rooted

## Lower bound on the kth root

$$
\begin{aligned}
\mathbb{E} \chi\left[\sum_{j=1}^{k} r_{j} r_{j}^{*}\right](x) & =\left(1-\frac{1}{n} \partial_{x}\right)^{k} x^{d} \\
& =x^{d-k}\left(1-\frac{1}{n} \partial_{x}\right)^{d} x^{k} \\
& =x^{d-k} L_{k}^{d-k}(n x)
\end{aligned}
$$

a scaled associated Laguerre polynomial

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\mathbb{E} \chi\left[\sum_{j=1}^{k} r_{j} r_{j}^{*}\right](x)=x^{d-k} L_{k}^{d-k}(n x)
$$

a scaled associated Laguerre polynomial

Let $R$ be a k-by-d matrix of independent $\mathcal{N}(0,1 / n)$

$$
\begin{aligned}
& \mathbb{E} \chi\left[R R^{T}\right]=L_{k}^{d-k}(n x) \\
& \mathbb{E} \chi\left[R^{T} R\right]=x^{d-k} L_{k}^{d-k}(n x)
\end{aligned}
$$

## Lower bound on the kth root

$$
\mathbb{E} \chi\left[\sum_{j=1}^{k} r_{j} r_{j}^{*}\right](x)=x^{d-k} L_{k}^{d-k}(n x)
$$

a scaled associated Laguerre polynomial

$$
\lambda_{k}\left(\mathbb{E} \chi\left[\sum r_{j} r_{j}^{*}\right](x)\right) \geq\left(1-\sqrt{\frac{k}{d}}\right)^{2} \frac{d}{n}
$$

(Krasikov '06)

## 3. With non-zero probability

$\lambda_{k}\left(\chi\left[\sum r_{j} r_{j}^{*}\right](x)\right) \geq \lambda_{k}\left(\mathbb{E} \chi\left[\sum r_{j} r_{j}^{*}\right](x)\right)$

Proof: the polynomials

$$
p_{i_{1}, i_{2}, \ldots, i_{k}}(x)=\chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{k}} v_{i_{k}}^{*}\right](x)
$$

form an interlacing family.

## Interlacing

Polynomial $p(x)=\prod_{i=1}^{d}\left(x-\alpha_{i}\right)$
interlaces $\quad q(x)=\prod_{i=1}^{d-1}\left(x-\beta_{i}\right)$
if $\quad \alpha_{1} \leq \beta_{1} \leq \alpha_{2} \leq \cdots \alpha_{d-1} \leq \beta_{d-1} \leq \alpha_{d}$

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if $\quad \alpha_{1} \leq \beta_{1} \leq \alpha_{2} \leq \cdots \alpha_{d-1} \leq \beta_{d-1} \leq \alpha_{d}$
For example, $p(x)$ interlaces $\partial_{x} p(x)$

## Interlacing

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if $\quad \alpha_{1} \leq \beta_{1} \leq \alpha_{2} \leq \cdots \alpha_{d-1} \leq \beta_{d-1} \leq \alpha_{d} \leq \beta_{d}$

If generalize to allow same degree,
Cauchy's interlacing theorem says
$\chi[A](x)$ interlaces $\chi\left[A+v v^{*}\right](x)$

## Common Interlacing

$p_{1}(x)$ and $p_{2}(x)$ have a common interlacing if can partition the line into intervals so that each contains one root from each polynomial

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If $p_{1}(x), p_{2}(x), \ldots, p_{n}(x)$ have a common interlacing
then $\min _{j} \lambda_{k}\left(p_{j}\right) \leq \lambda_{k}\left(\underset{j}{\mathbb{E}} p_{j}\right) \leq \max _{j} \lambda_{k}\left(p_{j}\right)$

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So, the average has a root between the smallest and largest kth roots

## Without a common interlacing



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## Interlacing Family of Polynomials

$\left\{p_{\sigma}(x)\right\}_{\sigma}$ is an interlacing family
if its members can be placed on the leaves of a tree so that when every node is labeled with the average of leaves below, siblings have common interlacings


## Interlacing Family of Polynomials

$\left\{p_{\sigma}(x)\right\}_{\sigma}$ is an interlacing family
For $\sigma \in\{1, . ., n\}^{k}$, set $p_{i_{1}, \ldots, i_{h}}=\mathbb{E}_{i_{h+1}, \ldots, i_{k}} p_{i_{1}, \ldots, i_{k}}$


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## Interlacing Family of Polynomials

Theorem:
There is an $i_{1}, \ldots, i_{k}$ so that

$$
\lambda_{k}\left(p_{i_{1}, \ldots, i_{k}}\right) \geq \lambda_{k}\left(p_{\emptyset}\right)
$$



## Interlacing Family of Polynomials

It remains to prove that the polynomials

$$
p_{i_{1}, i_{2}, \ldots, i_{k}}(x)=\chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{k}} v_{i_{k}}^{*}\right](x)
$$

form an interlacing family.
Will imply that with non-zero probability

$$
\lambda_{k}\left(\chi\left[\sum r_{j} r_{j}^{*}\right](x)\right) \geq \lambda_{k}\left(\mathbb{E} \chi\left[\sum r_{j} r_{j}^{*}\right](x)\right)
$$

## Common interlacings

$p_{1}(x), p_{2}(x), \ldots, p_{n}(x)$ have a common interlacing iff
for every convex combination $\sum_{j} \mu_{j}=1, \mu_{j} \geq 0$

$$
\sum_{j} \mu_{j} p_{j}(x)
$$

is real rooted


## Common interlacings

$p_{1}(x), p_{2}(x), \ldots, p_{n}(x)$ have a common interlacing iff
for every distribution $\mu$ on $\{1, \ldots, n\}$

$$
\mathbb{E}_{j \sim \mu} p_{j}(x)
$$

is real rooted


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$$
\mathbb{E}_{j \sim \mu} p_{j}(x)
$$

is real rooted

Proof: by similar picture.
(Dedieu '80, Fell '92, Chudnovsky-Seymour '07)

## An interlacing family

$$
p_{i_{1}, i_{2}, \ldots, i_{k}}(x)=\chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{k}} v_{i_{k}}^{*}\right](x)
$$

Nodes on the tree are labeled with

$$
p_{i_{1}, \ldots, i_{h}}(x)=\mathbb{E}_{i_{h+1}, \ldots, i_{k}} p_{i_{1}, \ldots, i_{k}}(x)
$$

We need to show that for each $i_{1}, \ldots, i_{h}$ the polynomials
$p_{i_{1}, \ldots, i_{h}, \underline{j}}(x)$ have a common interlacing

## An interlacing family

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p_{i_{1}, i_{2}, \ldots, i_{k}}(x)=\chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{k}} v_{i_{k}}^{*}\right](x)
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$$

We need to show that for each $i_{1}, \ldots, i_{h}$ the polynomials

$$
p_{i_{1}, \ldots, i_{h}, \underline{\underline{1}}}(x) \text { have a common interlacing }
$$

Proof:

$$
\begin{equation*}
=\left(1-\frac{1}{n} \partial_{x}\right)^{k-h-1} \chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}+v_{j} v_{j}^{*}\right] \tag{x}
\end{equation*}
$$

## An interlacing family

We need to show that for each $i_{1}, \ldots, i_{h}$ the polynomials
$p_{i_{1}, \ldots, i_{h}, \underline{j}}(x)$ have a common interlacing
Proof:

$$
=\left(1-\frac{1}{n} \partial_{x}\right)^{k-h-1} \chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}+v_{j} v_{j}^{*}\right](x)
$$

Cauchy's interlacing theorem implies

$$
\chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}+v_{j} v_{j}^{*}\right](x)
$$

All interlace

$$
\chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}\right](x)
$$

## An interlacing family

We need to show that for each $i_{1}, \ldots, i_{h}$ the polynomials
$p_{i_{1}, \ldots, i_{h}, j}(x)$ have a common interlacing
Proof:

$$
=\left(1-\frac{1}{n} \partial_{x}\right)^{k-h-1} \chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}+v_{j} v_{j}^{*}\right](x)
$$

Cauchy's interlacing theorem implies

$$
\chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}+v_{j} v_{j}^{*}\right](x)
$$

have a common interlacing

## An interlacing family

## Proof:

$$
=\left(1-\frac{1}{n} \partial_{x}\right)^{k-h-1} \chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}+v_{j} v_{j}^{*}\right](x)
$$

Cauchy's interlacing theorem implies

$$
\chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}+v_{j} v_{j}^{*}\right](x)
$$

have a common interlacing. So, for all distributions $\mu$

$$
\mathbb{E}_{j \sim \mu} \chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}+v_{j} v_{j}^{*}\right](x)
$$

is real rooted.

## An interlacing family

Proof:

$$
=\left(1-\frac{1}{n} \partial_{x}\right)^{k-h-1} \chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}+v_{j} v_{j}^{*}\right](x)
$$

Cauchy's interlacing theorem implies

$$
\chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}+v_{j} v_{j}^{*}\right](x)
$$

have a common interlacing. So, for all distributions $\mu$

$$
\mathbb{E}_{j \sim \mu} \chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}+v_{j} v_{j}^{*}\right](x)
$$

is real rooted. And,
$\mathbb{E}_{j \sim \mu}\left(1-\frac{1}{n} \partial_{x}\right)^{k-h-1} \chi\left[v_{i_{1}} v_{i_{1}}^{*}+\cdots+v_{i_{h}} v_{i_{h}}^{*}+v_{j} v_{j}^{*}\right](x)$ is also real rooted for every distribution $\mu$.

## Method of proof

Let $r_{1}, \ldots, r_{k}$ be chosen uniformly from $\left\{v_{1}, \ldots, v_{n}\right\}$

1. $\mathbb{E} \chi\left[\sum r_{j} r_{j}^{*}\right](x)$ is real rooted
2. $\lambda_{k}\left(\mathbb{E} \chi\left[\sum r_{j} r_{j}^{*}\right](x)\right) \geq\left(1-\sqrt{\frac{k}{d}}\right)^{2} \frac{d}{n}$
3. With non-zero probability

$$
\lambda_{k}\left(\chi\left[\sum r_{j} r_{j}^{*}\right](x)\right) \geq \lambda_{k}\left(\mathbb{E} \chi\left[\sum r_{j} r_{j}^{*}\right](x)\right)
$$

Because is an interlacing family of polynomials

## In part 2

Will use the same approach to prove
Weaver's conjecture, and thereby Kadison-Singer
But, employ multivariate analogs of these arguments
and a direct bound on the roots of the polynomials.

## Main Theorem

For all $\alpha>0$
if all $\left\|v_{i}\right\| \leq \alpha$ and $\sum v_{i} v_{i}^{*}=I$
then exists a partition into $S_{1}$ and $S_{2}$ with

$$
\operatorname{eigs}\left(\sum_{i \in S_{j}} v_{i} v_{i}^{*}\right) \leq \frac{1}{2}+3 \alpha
$$

Implies Akemann-Anderson Paving Conjecture, which implies Kadison-Singer

## We want



## We want

$$
\operatorname{roots}\left(\chi\left[\begin{array}{cc}
\sum_{i \in S_{1}} v_{i} v_{i}^{*} & 0 \\
0 & \sum_{i \in S_{2}} v_{i} v_{i}^{*}
\end{array}\right](x)\right) \leq \frac{1}{2}+3 \alpha
$$

## We want

$$
\operatorname{roots}\left(\chi\left[\begin{array}{cc}
\sum_{i \in S_{1}} v_{i} v_{i}^{*} & 0 \\
0 & \sum_{i \in S_{2}} v_{i} v_{i}^{*}
\end{array}\right](x)\right) \leq \frac{1}{2}+3 \alpha
$$

Consider expected polynomial with a random partition.

## Method of proof

1. Prove expected characteristic polynomial has real roots
2. Prove its largest root is at most $1 / 2+3 \alpha$
3. Prove is an interlacing family, so exists a partition whose polynomial has largest root at most $1 / 2+3 \alpha$

## The Expected Polynomial

Indicate choices by $\sigma_{1}, \ldots, \sigma_{n}: i \in S_{\sigma_{i}}$

$$
p_{\sigma_{1}, \ldots, \sigma_{n}}(x)=\chi\left[\begin{array}{cc}
\sum_{i: \sigma_{i}=1} v_{i} v_{i}^{*} & 0 \\
0 & \sum_{i: \sigma_{i}=2} v_{i} v_{i}^{*}
\end{array}\right]
$$

$(x)$

## The Expected Polynomial

$$
a_{i}=\binom{v_{i}}{0} \text { for } i \in S_{1} \quad a_{i}=\binom{0}{v_{i}} \text { for } i \in S_{2}
$$

$$
\left(\begin{array}{cc}
\sum_{i \in S_{1}} v_{i} v_{i}^{*} & 0 \\
0 & \sum_{i \in S_{2}} v_{i} v_{i}^{*}
\end{array}\right)=\sum_{i} a_{i} a_{i}^{*}
$$

## Mixed Characteristic Polynomials

For $a_{1}, \ldots, a_{n}$ independently chosen random vectors

$$
\mathbb{E} \chi\left[\sum_{i} a_{i} a_{i}^{*}\right]
$$

is their mixed characteristic polynomial.

Theorem: It only depends on $A_{i}=\mathbb{E} a_{i} a_{i}^{*}$ and, is real-rooted

## Mixed Characteristic Polynomials

For $a_{1}, \ldots, a_{n}$ independently chosen random vectors

$$
\mathbb{E} \chi\left[\sum_{i} a_{i} a_{i}^{*}\right]=\mu\left(A_{1}, \ldots, A_{n}\right)
$$

is their mixed characteristic polynomial.

Theorem: It only depends on $A_{i}=\mathbb{E} a_{i} a_{i}^{*}$ and, is real-rooted

$$
\operatorname{Tr}\left(A_{i}\right)=\operatorname{Tr}\left(\mathbb{E} a_{i} a_{i}^{*}\right)=\mathbb{E} \operatorname{Tr}\left(a_{i} a_{i}^{*}\right)=\mathbb{E}\left\|a_{i}\right\|^{2}
$$

## Mixed Characteristic Polynomials

For $a_{1}, \ldots, a_{n}$ independently chosen random vectors

$$
\mathbb{E} \chi\left[\sum_{i} a_{i} a_{i}^{*}\right]=\mu\left(A_{1}, \ldots, A_{n}\right)
$$

is their mixed characteristic polynomial.

The constant term is the mixed discriminant of

$$
A_{1}, \ldots, A_{n}
$$

## The constant term

When diagonal and $d=n, c_{d}$ is a matrix permanent.


## The constant term

When diagonal and $d=n, c_{d}$ is a matrix permanent. Van der Waerden's Conjecture becomes

$$
\text { If } \sum A_{i}=I \text { and } \operatorname{Tr}\left(A_{i}\right)=1
$$

$c_{d}$ is minimized when $A_{i}=\frac{1}{n} I$

Proved by Egorychev and Falikman '81.
Simpler proof by Gurvits (see Laurent-Schrijver)

## The constant term

For Hermitian matrices, $c_{d}$ is the mixed discriminant Gurvits proved a lower bound on $c_{d}$ :

$$
\begin{aligned}
& \text { If } \sum A_{i}=I \text { and } \operatorname{Tr}\left(A_{i}\right)=1 \\
& \qquad c_{d} \text { is minimized when } A_{i}=\frac{1}{n} I
\end{aligned}
$$

This was a conjecture of Bapat.

## Other coefficients

One can generalize Gurvits's results to prove lower bounds on all the coefficients.

$$
\begin{aligned}
& \text { If } \sum A_{i}=I \text { and } \operatorname{Tr}\left(A_{i}\right)=1 \\
& \qquad\left|c_{k}\right| \text { is minimized when } A_{i}=\frac{1}{n} I
\end{aligned}
$$

But, this does not imply useful bounds on the roots

## Real Stable Polynomials

A multivariate generalization of real rootedness.

Complex roots of $p \in \mathbb{R}[z]$ come in conjugate pairs.

So, real rooted iff no roots with positive complex part.

## Real Stable Polynomials

$p \in \mathbb{R}\left[z_{1}, \ldots, z_{n}\right]$
is real stable if $\operatorname{imag}\left(z_{i}\right)>0$ for all $i$ implies $p\left(z_{1}, \ldots, z_{n}\right) \neq 0$
it has no roots in the upper half-plane
Isomorphic to Gårding's hyperbolic polynomials
Used by Gurvits (in his second proof)

## Real Stable Polynomials

$p \in \mathbb{R}\left[z_{1}, \ldots, z_{n}\right]$
is real stable if $\operatorname{imag}\left(z_{i}\right)>0$ for all $i$ implies $p\left(z_{1}, \ldots, z_{n}\right) \neq 0$
it has no roots in the upper half-plane
Isomorphic to Gårding's hyperbolic polynomials
Used by Gurvits (in his second proof)

See surveys of Pemantle and Wagner

## Real Stable Polynomials

Borcea-Brändén ‘08:
For PSD matrices $A_{1}, \ldots, A_{n}$

$$
\operatorname{det}\left[z_{1} A_{1}+\cdots+z_{n} A_{n}\right]
$$

is real stable

## Real Stable Polynomials

$p\left(z_{1}, \ldots, z_{n}\right)$ real stable
implies $\left(1-\partial_{z_{i}}\right) p\left(z_{1}, \ldots, z_{n}\right)$ is real stable
(Lieb Sokal '81)
$p\left(z_{1}, \ldots, z_{n}\right)$ real stable implies $p(x, x, \ldots, x)$ is real rooted

## Real Roots

$\mu\left(A_{1}, \ldots, A_{n}\right)(x)=$

$$
\left.\left(\prod_{i=1}^{n} 1-\partial_{z_{i}}\right) \operatorname{det}\left(\sum_{i=1}^{n} z_{i} A_{i}\right)\right|_{z_{1}=\cdots=z_{n}=x}
$$

So, every mixed characteristic polynomial is real rooted.

## Our Interlacing Family

Indicate choices by $\sigma_{1}, \ldots, \sigma_{n}: i \in S_{\sigma_{i}}$

$$
\begin{align*}
& p_{\sigma_{1}, \ldots, \sigma_{n}}(x)=\chi\left[\begin{array}{cc}
\sum_{i: \sigma_{i}=1} v_{i} v_{i}^{*} & 0 \\
0 & \sum_{i: \sigma_{i}=2} v_{i} v_{i}^{*}
\end{array}\right]  \tag{x}\\
& p_{\sigma_{1}, \ldots, \sigma_{k}}(x)=\underset{\sigma_{k+1}, \ldots, \sigma_{n}}{\mathbb{E}}\left[p_{\sigma_{1}, \ldots, \sigma_{n}}\right](x)
\end{align*}
$$

## Interlacing

$p_{1}(x)$ and $p_{2}(x)$ have a common interlacing iff
$\lambda p_{1}(x)+(1-\lambda) p_{2}(x)$ is real rooted for all $0 \leq \lambda \leq 1$
We need to show that

$$
\lambda p_{\sigma_{1}, \ldots, \sigma_{k}, 1}(x)+(1-\lambda) p_{\sigma_{1}, \ldots \sigma_{k}, 2}(x)
$$

is real rooted.

## Interlacing

$p_{1}(x)$ and $p_{2}(x)$ have a common interlacing iff
$\lambda p_{1}(x)+(1-\lambda) p_{2}(x)$ is real rooted for all $0 \leq \lambda \leq 1$
We need to show that

$$
\lambda p_{\sigma_{1}, \ldots, \sigma_{k}, 1}(x)+(1-\lambda) p_{\sigma_{1}, \ldots \sigma_{k}, 2}(x)
$$

is real rooted.
It is a mixed characteristic polynomial, so is real-rooted.
Set $\sigma_{k+1}=1$ with probability $\lambda$
Keep $\sigma_{i}$ uniform for $i>k+1$

## An upper bound on the roots

Theorem: If $\sum A_{i}=I$ and $\operatorname{Tr}\left(A_{i}\right) \leq \epsilon$ then

$$
\max -\operatorname{root}\left(\mu\left(A_{1}, \ldots, A_{n}\right)(x)\right) \leq(1+\sqrt{\epsilon})^{2}
$$

## An upper bound on the roots

Theorem: If $\sum A_{i}=I$ and $\operatorname{Tr}\left(A_{i}\right) \leq \epsilon$ then $\max -\operatorname{root}\left(\mu\left(A_{1}, \ldots, A_{n}\right)(x)\right) \leq(1+\sqrt{\epsilon})^{2}$

An upper bound of 2 is trivial (in our special case).
Need any constant strictly less than 2.

## An upper bound on the roots

Theorem: If $\sum A_{i}=I$ and $\operatorname{Tr}\left(A_{i}\right) \leq \epsilon$ then $\max -\operatorname{root}\left(\mu\left(A_{1}, \ldots, A_{n}\right)(x)\right) \leq(1+\sqrt{\epsilon})^{2}$
$\mu\left(A_{1}, \ldots, A_{n}\right)(x)=$

$$
\left.\left(\prod_{i=1}^{n} 1-\partial_{z_{i}}\right) \operatorname{det}\left(\sum_{i=1}^{n} z_{i} A_{i}\right)\right|_{z_{1}=\cdots=z_{n}=x}
$$

## An upper bound on the roots

Define: $\left(w_{1}, \ldots, w_{n}\right)$ is an upper bound on the roots of $p\left(z_{1}, \ldots, z_{n}\right)$ if

$$
p\left(z_{1}, \ldots, z_{n}\right)>0 \text { for }\left(z_{1}, \ldots, z_{n}\right) \geq\left(w_{1}, \ldots, w_{n}\right)
$$



## An upper bound on the roots

Define: $\left(w_{1}, \ldots, w_{n}\right)$ is an upper bound on the roots of $p\left(z_{1}, \ldots, z_{n}\right)$ if

$$
p\left(z_{1}, \ldots, z_{n}\right)>0 \text { for }\left(z_{1}, \ldots, z_{n}\right) \geq\left(w_{1}, \ldots, w_{n}\right)
$$



Eventually set

$$
z_{1}, \ldots, z_{n}=x
$$

so want

$$
w_{1}=\cdots=w_{n}
$$

## Action of the operators



## Action of the operators



## Action of the operators



## The roots of $\left(1-\partial_{x}\right) p(x)$

Define:

$$
\begin{aligned}
& \alpha-\max \left(\lambda_{1}, \ldots, \lambda_{n}\right)=\max \left\{u: \sum \frac{1}{u-\lambda_{i}}=\alpha\right\} \\
& \alpha-\max (p(x))=\alpha-\max (\operatorname{roots}(p))
\end{aligned}
$$

Theorem (Batson-S-Srivastava):
If $p(x)$ is real rooted and $\alpha>0$

$$
\alpha-\max \left(\left(1-\partial_{x}\right) p(x)\right) \leq \alpha-\max (p(x))+\frac{1}{1-\alpha}
$$

## The roots of $\left(1-\partial_{x}\right) p(x)$

Theorem (Batson-S-Srivastava):
If $p(x)$ is real rooted and $\alpha>0$
$\alpha-\max \left(\left(1-\partial_{x}\right) p(x)\right) \leq \alpha-\max (p(x))+\frac{1}{1-\alpha}$
Proof: Define $\Phi_{p}(u)=\frac{p^{\prime}(u)}{p(u)}=\sum_{i} \frac{1}{u-\lambda_{i}}=\partial_{u} \log p(u)$
Set $u=\alpha-\max (p(x))$, so $\Phi_{p}(u)=\alpha$

Suffices to show for all $\delta \geq \frac{1}{1-\alpha}$

$$
\Phi_{p-p^{\prime}}(u+\delta) \leq \Phi_{p}(u)
$$

## The roots of $\left(1-\partial_{x}\right) p(x)$

Define $\Phi_{p}(u)=\frac{p^{\prime}(u)}{p(u)}=\sum_{i} \frac{1}{u-\lambda_{i}}$
Set $u=\alpha-\max (p(x))$, so $\Phi_{p}(u)=\alpha$
Suffices to show for all $\delta \geq \frac{1}{1-\alpha}$

$$
\begin{gathered}
\Phi_{p-p^{\prime}}(u+\delta) \leq \Phi_{p}(u) \\
\\
\downarrow \text { (algebra) } \\
\Phi_{p}(u)-\Phi_{p}(u+\delta) \geq \frac{-\Phi_{p}^{\prime}(u+\delta)}{1-\Phi_{p}(u+\delta)}
\end{gathered}
$$

## The roots of $\left(1-\partial_{x}\right) p(x)$

Define $\Phi_{p}(u)=\frac{p^{\prime}(u)}{p(u)}=\sum_{i} \frac{1}{u-\lambda_{i}}$

$$
\Phi_{p}(u)-\Phi_{p}(u+\delta) \geq \frac{-\Phi_{p}^{\prime}(u+\delta)}{1-\Phi_{p}(u+\delta)}
$$

$\Phi_{p}(u)$ convex for $u>\max (p(x))$ implies

$$
\Phi_{p}(u)-\Phi_{p}(u+\delta) \geq \delta\left(-\Phi_{p}^{\prime}(u+\delta)\right)
$$

Monotone decreasing implies only need

$$
\delta \geq \frac{1}{1-\Phi_{p}(u+\delta)}
$$

## The roots of $\left(1-\partial_{x}\right) p(x)$

$\Phi_{p}(u)$ convex for $u>\max (p(x))$ implies

$$
\Phi_{p}(u)-\Phi_{p}(u+\delta) \geq \delta\left(-\Phi_{p}^{\prime}(u+\delta)\right)
$$

Monotone decreasing implies only need

$$
\delta \geq \frac{1}{1-\Phi_{p}(u+\delta)}
$$

and that $\quad \delta \geq \frac{1}{1-\alpha}=\frac{1}{1-\Phi_{p}(u)} \quad$ suffices.

## The roots of $\left(1-\partial_{x}\right) p(x)$

$\Phi_{p}(u)$ convex for $u>\max (p(x))$ implies

$$
\Phi_{p}(u)-\Phi_{p}(u+\delta) \geq \delta\left(-\Phi_{p}^{\prime}(u+\delta)\right)
$$

Monotone decreasing_implies only need

$$
\delta \geq \frac{1}{1-\Phi_{p}(u+\delta)}
$$

and that $\quad \delta \geq \frac{1}{1-\alpha}=\frac{1}{1-\Phi_{p}(u)} \quad$ suffices.

## The roots of $\left(1-\partial_{x}\right) p(x)$

Theorem (Batson-S-Srivastava):
If $p(x)$ is real rooted and $\alpha>0$

$$
\alpha-\max \left(\left(1-\partial_{x}\right) p(x)\right) \leq \alpha-\max (p(x))+\frac{1}{1-\alpha}
$$

Gives a sharp upper bound on the roots of associated Laguerre polynomials.

The analogous argument with the min gives the lower bound that we claimed before.

## An upper bound on the roots

Theorem: If $\sum A_{i}=I$ and $\operatorname{Tr}\left(A_{i}\right) \leq \epsilon$ then $\max -\operatorname{root}\left(\mu\left(A_{1}, \ldots, A_{n}\right)(x)\right) \leq(1+\sqrt{\epsilon})^{2}$
$\mu\left(A_{1}, \ldots, A_{n}\right)(x)=$

$$
\left.\left(\prod_{i=1}^{n} 1-\partial_{z_{i}}\right) \operatorname{det}\left(\sum_{i=1}^{n} z_{i} A_{i}\right)\right|_{z_{1}=\cdots=z_{n}=x}
$$

## A robust upper bound

Define: $\left(w_{1}, \ldots, w_{n}\right)$ is an $\alpha$-upper bound on $p\left(z_{1}, \ldots, z_{n}\right)$
if it is an $\alpha_{i}$-max, in each $z_{i}$, and $\alpha_{i} \leq \alpha$

Theorem:
If $w$ is an $\alpha$-upper bound on $p$, then $w+\delta e_{j}$ is an $\alpha$-upper bound on $p-\partial_{z_{j}} p$,

$$
\text { for } \delta \geq \frac{1}{1-\alpha}
$$

## A robust upper bound

Theorem:
If $w$ is an $\alpha$-upper bound on $p$, and $\delta \geq \frac{1}{1-\alpha}$ $w+\delta e_{j}$ is an $\alpha$-upper bound on $p-\partial_{z_{j}} p$,

Proof:
Same as before, but need to know that

$$
\frac{\partial_{z_{i}} p\left(z_{1}, \ldots, z_{n}\right)}{p\left(z_{1}, \ldots, z_{n}\right)}
$$

is decreasing and convex in $z_{j}$, above the roots

## A robust upper bound

## Proof:

Same as before, but need to know that

$$
\frac{\partial_{z_{i}} p\left(z_{1}, \ldots, z_{n}\right)}{p\left(z_{1}, \ldots, z_{n}\right)}
$$

is decreasing and convex in $z_{j}$, above the roots
Follows from a theorem of Helton and Vinnikov '07:
Every bivariate real stable polynomial can be written

$$
\operatorname{det}(A+B x+C y)
$$

## A robust upper bound

## Proof:

Same as before, but need to know that

$$
\frac{\partial_{z_{i}} p\left(z_{1}, \ldots, z_{n}\right)}{p\left(z_{1}, \ldots, z_{n}\right)}
$$

is decreasing and convex in $z_{j}$, above the roots
Or, as pointed out by Renegar,
from a theorem Bauschke, Güler, Lewis, and Sendov '01
Or, by a theorem of Brändén ' 07 .
Or, see Terry Tao's blog for a (mostly) self-contained proof

## An upper bound on the roots

Theorem: If $\sum A_{i}=I$ and $\operatorname{Tr}\left(A_{i}\right) \leq \epsilon$ then

$$
\max -\operatorname{root}\left(\mu\left(A_{1}, \ldots, A_{n}\right)(x)\right) \leq(1+\sqrt{\epsilon})^{2}
$$

## A probabilistic interpretation

For $a_{1}, \ldots, a_{n}$ independently chosen random vectors with finite support
such that $\mathbb{E}\left[\sum_{i} a_{i} a_{i}^{T}\right]=I$ and $\left\|\mathbb{E}\left[a_{i} a_{i}^{T}\right]\right\| \leq \epsilon$
then $\operatorname{Pr}\left[\left\|\sum_{i} a_{i} a_{i}^{T}\right\| \leq(1+\sqrt{\epsilon})^{2}\right]>0$

## Main Theorem

For all $\alpha>0$
if all $\left\|v_{i}\right\| \leq \alpha$
then exists a partition into $S_{1}$ and $S_{2}$ with

$$
\operatorname{eigs}\left(\sum_{i \in S_{j}} v_{i} v_{i}^{T}\right) \leq \frac{1}{2}+3 \alpha
$$

Implies Akemann-Anderson Paving Conjecture, which implies Kadison-Singer

## Anderson's Paving Conjecture ‘79

Reduction by Casazza-Edidin-Kalra-Paulsen '07 and Harvey '13:
There exist an $\epsilon>0$ and a $k$ so that
if all $\left\|v_{i}\right\|^{2} \leq 1 / 2$ and $\sum v_{i} v_{i}^{T}=I$
then exists a partition of $\{1, \ldots, n\}$ into $k$ parts s.t.

$$
\operatorname{eigs}\left(\sum_{i \in S_{j}} v_{i} v_{i}^{T}\right) \leq 1-\epsilon
$$

Can prove using the same technique

## A conjecture

If $\sum A_{i}=I$ and $\operatorname{Tr}\left(A_{i}\right) \leq \epsilon$ then max-root $\left(\mu\left(A_{1}, \ldots, A_{n}\right)(x)\right)$
is largest when $A_{i}=\frac{\epsilon}{d} I$

## Questions

Can the partition be found in polynomial time?

What else can one construct this way?

How do operations that preserve real rootedness move the roots and the Stieltjes transform?

