

Linear Programming & Convex Sets.

An LP has variables $x \in \mathbb{R}^d$,

an objective function $\max c^T x$, specified by $c \in \mathbb{R}^d$

subject to m constraints $a_i^T x \leq b_i$,

where $a_i \in \mathbb{R}^d$, $b_i \in \mathbb{R}$, for $1 \leq i \leq m$.

We typically collect the a_i into an m -by- d matrix, A ,
and the b_i into an m -vector, b , and write

$Ax \leq b$, where this means the inequality holds
in every coordinate:

$x \leq y$ iff for all i $x(i) \leq y(i)$.

We will see many forms of linear programs.

If want $\bar{a}^T x \geq \bar{b}$, write it as $(-\bar{a})^T x \leq -\bar{b}$.

If want $\bar{a}^T x = \bar{b}$, write it as $\bar{a}^T x \leq \bar{b}$
and $(-\bar{a})^T x \leq -\bar{b}$

The set $\{x: Ax \leq b\}$ is a polyhedron.

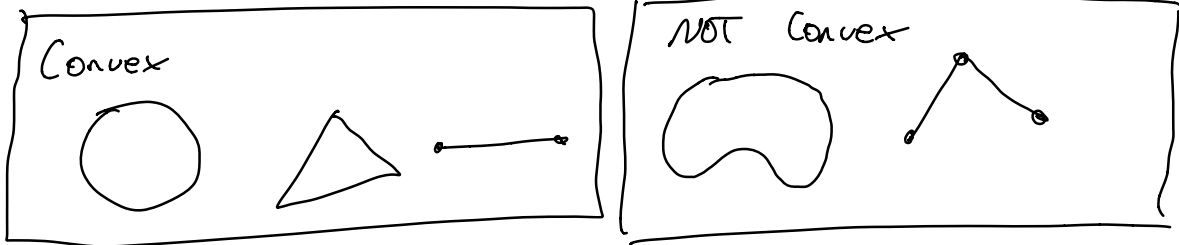
In particular, it is convex.

We will spend most of this lecture understanding
convex sets.

Def $C \subseteq \mathbb{R}^n$ is convex if for all $x, y \in C$

and all $0 \leq t \leq 1$, $tx + (1-t)y \in C$.

That is, the line segment from x to y is in C .



Examples

A subspace: $\{x: Ax = 0\}$

$$Ax = 0 \text{ and } Ay = 0 \Rightarrow A(tx + (1-t)y) = 0.$$

An affine space: $\{x: Ax = b\}$

$$Ax = b \text{ and } Ay = b \Rightarrow$$

$$A(tx + (1-t)y) = tAx + (1-t)Ay = tb + (1-t)b = b.$$

A line segment: $\{x: a \leq x \leq b\}$, $a, b \in \mathbb{R}$

An open line segment: $\{x: a < x < b\}$, $a, b \in \mathbb{R}$

Positive reals: $x > 0$

A halfspace: $x \in \mathbb{R}^d$ s.t. $a^T x \leq b$

$$a^T (tx + (1-t)y) = ta^T x + (1-t)a^T y$$

$$\leq tb + (1-t)b = b$$

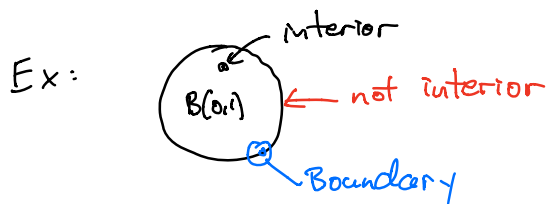
An open halfspace: $x \in \mathbb{R}^d$ s.t. $a^T x < b$

A norm ball: $\{x: \|x\| \leq 1\}$, for any norm $\|\cdot\|$

Let $B(x, r)$ be the ball of radius r around x
 $= \{y: \|x-y\| \leq r\}$

Some topology

For $S \subseteq \mathbb{R}^d$, x is in the interior of S
if $\exists \varepsilon > 0$ s.t. $B(x, \varepsilon) \subseteq S$



x is on the boundary of S if $\forall \varepsilon > 0$,
 $B(x, \varepsilon)$ contains points in S and not in S

If $S = \{x: 0 < x \leq 1\}$, The boundary = $\{0, 1\}$
 $1 \in S$ but $0 \notin S$

S is closed if $\text{boundary}(S) \subseteq S$

S is open if $S = \text{interior}(S)$

e.g. $\{x: 0 < x < 1\}$ $\forall \varepsilon < \frac{1}{2}$, $B(x, \varepsilon) \subseteq S$

S is bounded if $\exists r$ s.t. $S \subseteq B(0, r)$.

A halfspace is not bounded

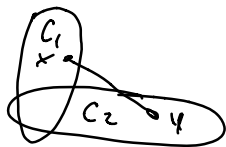
$S \subseteq \mathbb{R}^d$ is compact if it is closed and bounded.
Might say what it means later

Intersections: if C_1 and C_2 are convex,
so is $C_1 \cap C_2$

The empty set is convex,
as is a singleton.

A polyhedron: $x: Ax \leq b$.

Unions DO NOT preserve convexity



$C_1 \cup C_2$ not necessarily
convex

Given $x_1, \dots, x_k \in \mathbb{R}^d$, a convex combination
of x_1, \dots, x_k is a point $x = \sum_i t_i x_i$

where $\sum_i t_i = 1$ and $t_i \geq 0$.

The convex hull, written $\text{CH}(x_1, \dots, x_k)$ is the
set of all convex combinations

$$\left\{ \sum_i t_i x_i : t_i \geq 0, \mathbf{1}^T t = 1 \right\}$$

This is always convex.

proof If $y = \sum r_i x_i$ $r \geq 0$ $\mathbf{1}^T r = 1$

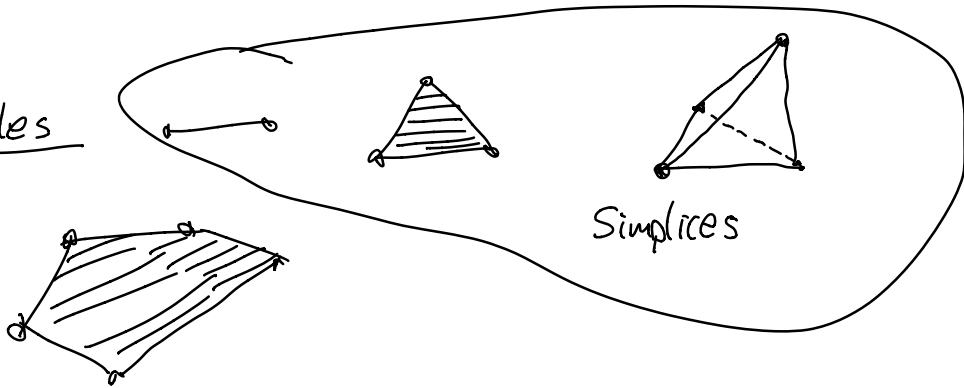
$z = \sum s_i x_i$ $s \geq 0$ $\mathbf{1}^T s = 1$

and $t \in [0, 1]$

$t y + (1-t) z = \sum u_i x_i$

with $u_i = t r_i + (1-t) s_i$

Examples



We say x_0, \dots, x_k are affinely independent

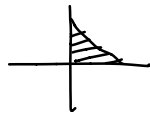
if $x_1 - x_0, \dots, x_k - x_0$ are linearly independent.

is equivalent to $\begin{pmatrix} x_0 \\ 1 \end{pmatrix}, \dots, \begin{pmatrix} x_k \\ 1 \end{pmatrix}$ being independent.

If x_0, \dots, x_k are affinely independent

then $\text{Conv}(x_0, \dots, x_k)$ is a simplex.

The standard simplex in \mathbb{R}^d is x s.t. $x \geq 0$
 $\mathbf{1}^T x = 1$



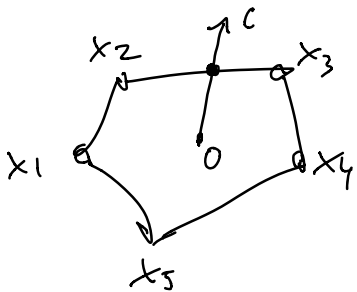
The probability simplex in \mathbb{R}^d is x s.t. $x \geq 0$
 $\mathbf{1}^T x = 1$.

Dan's Favorite LP:

given $x_1, \dots, x_m \in \mathbb{R}^d$ s.t. $0 \in \text{CH}(x_1, \dots, x_m)$

and $c \in \mathbb{R}^d$, max α

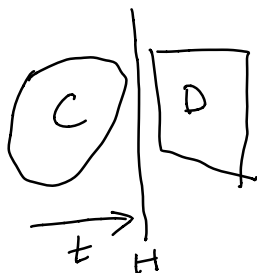
s.t. $\alpha c \in \text{CH}(x_1, \dots, x_m)$



Separating Hyperplane Thm 1

If C and D are disjoint closed convex sets and C is compact, then exists a hyperplane that separates them.

That is, $\exists t$ s.t. $\forall x \in C$ and $y \in D$ $t^T x < t^T y$.



if not convex.



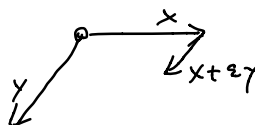
Fact If f is continuous and C is compact,

$\exists x \in C$ s.t. $f(y) \geq f(x) \forall y \in C$.

"achieves its minimum"

Lem If $x^T y < 0$ then for all $0 < \varepsilon < \frac{\|y\|^2}{-x^T y}$

$$\|x + \varepsilon y\|_2^2 < \|x\|_2^2.$$



proof $\|x + \varepsilon y\|_2^2 = \|x\|^2 + \varepsilon x^T y + \varepsilon^2 \|y\|^2$

This is $< \|x\|_2^2$ if $\varepsilon x^T y + \varepsilon^2 \|y\|^2 < 0$

$$\Leftrightarrow x^T y + \varepsilon \|y\|^2 < 0 \quad (\varepsilon > 0)$$

$$\Leftrightarrow \varepsilon < \|y\|^2 / (-x^T y)$$

proof of Thm 1

Let $c \in C$ and $d \in D$ minimize $\|c - d\|$

$\exists c \in C$ that minimizes $\text{dist}(C, D)$ because C compact.

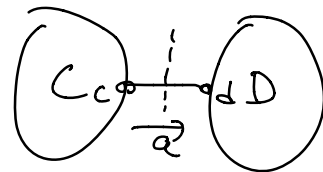
To prove is a d , look in $d \cap B(c, \text{dist}(c, D))$,
which is also compact.

As C and D are disjoint, $\|c - d\| \neq 0$.

Let $a = d - c$.

$$a^T c = d^T c - \|c\|_2^2$$

$$a^T d = \|d\|_2^2 - d^T c$$



Claim $\forall x \in C, a^T x \leq a^T c$ and $\forall x \in D, a^T x \geq a^T d$

$a^T c \neq a^T d$ because $a^T d - a^T c = \|d - c\|_2^2 > 0$.

proof. let $x \in C$. So, $\forall t \in [0, 1]$ $tx + (1-t)c \in C$
 $c + t(x-c)$

As $\text{dist}(c + t(x-c), d) \geq \text{dist}(c, d)$ for all $t \in [0, 1]$,
 $\|c - d + t(x - c)\| \geq \|c - d\|$

$$\text{lem 1} \Rightarrow (x - c)^T (c - d) \geq 0$$

$$\Rightarrow (x - c)^T a \leq 0$$

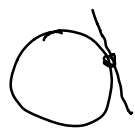
$$\Rightarrow x^T a \leq c^T a$$

The case of D is similar.

Supporting Hyperplane Theorem:

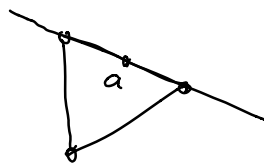
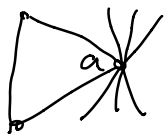
For all convex C and $a \in \text{boundary}(C)$,

$\exists t \neq 0$ such that $\forall x \in C, t^T x \leq t^T a$.



hyperplane through a
with C on one side.

$H = \{x: t^T x = t^T a\}$ is the supporting hyperplane at a



Proof idea :

For each $\varepsilon > 0$, let $a_\varepsilon \in B(a, \varepsilon)$ but $a_\varepsilon \notin C$.
Is a hyperplane that separates a_ε from C .
Take a limit of these.

If there is time, one last convex set:

Positive Semidefinite Matrices:

$A \in S_+^n$ if A is $n \times n$, symmetric, and
for all x $x^T A x \geq 0$.

If $A \in S_+^n$, can prove it by writing $A = LL^T$ —
the Cholesky Factorization.

If A is symmetric, but $A \notin S_+^n$,

let $A = U \Lambda U^T$ be spectral decomposition and $\lambda_n < 0$.

A hyperplane separating A from S_+^n is given by

$$\{ \text{symmetric } X : v_n^T X v_n = 0 \}$$

$$v_n^T X v_n \geq 0 \text{ for } X \in S_+^n. \quad v_n^T A v_n = \lambda_n < 0,$$

And is a hyperplane because

$$v_n^T X v_n = \sum_{1 \leq i, j \leq n} X(i, j) v(i) v(j) \text{ is linear in } X.$$