$$\frac{Convex Functions}{A \text{ function } f^{-}[R^{n} \rightarrow [R \text{ is } convex \text{ if for all}]}$$

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$$x \text{ and } y \text{ and all } O \leq t \leq 1,$$

$$f(tx + (l-t)y) \leq tf(x) + (l-t)f(y)$$

$$(X)$$



- f is  $\frac{\text{strictl}_{1}}{f(t+x+(1-t))} < tf(x) + (1-t)f(x)$
- |1×112 is strictly convex, but 11×1100 and 11×11, are merely convex.

Example: 
$$f(x) = \frac{1}{x^2}$$
  
is convex if don $[f] = (0, 00]$ ,  
bat not on  $\mathbb{R}$   
Same for  $x^3$   
If f is convex but don $[f] \neq \mathbb{R}^n$ ,  
the extension of f, f, is convex on  $\mathbb{R}^n$ , where  
 $F(x) = \int f(x) + f(x) + f(x) + f(x)$ 

This is one of the broadest classes of problems that we can solve efficiently.

If f is strictly convex, the minimum is unique: if X1, X2 & arg min f(X), but X1 + X2, then f( ± X1 + ± X2) < ± f(X1) + f(X2), contradicting minimality of X1, X2.

Local minima of convex functions are global minima.  

$$x$$
 is a local minimum of  $f$  if  $\exists z > 0$  sit.  
 $\||x-y\|| \in z \implies f(x) \in f(y)$ .  
 $x$  is a global minimum if  $\forall y = f(x) = f(x)$ .

If x were a local minimum but NOT global,  
then I y st. f(x) < f(x).  
For every E>O there is a te[o,1] s.t.  
$$\|x - ((1-t)x+t+1)\| \in E$$
.  
For this t, f((1-t)x+t+1) < (1-t) f(x) + tf(x) < f(x).  
This contradicts the assortion that x is a  
local minimum.

Examples of convex functions.  
One variable:  

$$e^{ax}$$
 for a e IR  
 $x^{d}$  dzl or  $x \in 0$ ,  $x \in [0, \infty]$   
 $-x^{d}$  O  $\in x \in [1, x \in [0, \infty]]$   
 $-\log(x)$ ,  $x \in [0, \infty]$ , because  $\log(x)$  is concave

Vectors

Affine functions : 
$$f(x) = a^T x + b$$

Maximum: Hie maximum of convex is convex  $f(x) = max(f_1(x), f_2(x))$ 

Ex. maximum of linear is convex



$$E_{\times}$$
, max (x(i), x(2),..., x(u))

Affine composition:  
if g is convex, then so is 
$$f(x) = g(Ax+b)$$

Examples: 
$$\|\|_{2}$$
 is convex  $\Rightarrow \|Ax\|_{2}$  convex  
 $\Rightarrow x^{T}(A^{T}A) \times is$  convex.  
Every Positive Semidefinite  $Q = A^{T}A$ , for some A.  
So,  $x^{T}Q \times is$  convex.

Least squares:  $||Ax-b||_2^2$  is convex  $l_1 - regularized : ||Ax-b||_2 + \lambda ||x||_1$  is convex  $logistic loss : \sum_i log(1 + e^{Y_i(aIx+b)}) = f(x)$ 

$$g(z) = log(l + e^{\gamma z})$$
 is convexing, because  
 $g''(z) = \frac{\gamma^2 e^{\gamma z}}{(l + e^{\gamma z})^2}$   
 $Q_{z}^{T}x + b$  is affine.

Composition rules (one of many)  
If 
$$f(x) = h(g(x))$$
 where  $h: |\mathbb{R} \to \mathbb{R}, g: |\mathbb{R}^n \to \mathbb{R},$   
g convex & h convex, non-decreasing => f convex







f is a convex function iff its epigoph is a convex set.

In one variable, a twice-differentiable function f(x) is convex iff f"(x) = 0, Hx & dom(f). This says that the function always lies above tonsent lines, like

Thm I  
In 
$$\Pi^n$$
, if f is differentiable then f is convex iff  
 $\forall x, y \in dom(f)$ ,  $f(y) \ge f(x) + \nabla f(x)^T(y-x)$  (\*)

Note that 
$$h(1) \triangleq f(1) + Vf(x)^{(\gamma-x)}$$
 is the  
linear function in  $\gamma$  such that  $h(x) = f(x)$ ,  
and is a supporting hyperplane of epigraph  
at  $(x, f(x))$ 

$$\frac{\text{proof of Hum}|}{\text{We first prove } (*)} \Rightarrow \text{convex}$$

$$\text{Let } x, y \in \text{dom}(f), \quad 0 < \lambda < l, \quad z = \lambda x + (l \cdot \lambda) y.$$

$$\text{Let } g = \mathcal{D}f(z).$$

$$(*) \Rightarrow \lambda f(x) \geq \lambda f(z) + \lambda g^{T}(x - z)$$

$$(l \cdot \lambda) f(z) \geq (l \cdot \lambda) f(z) + (l \cdot \lambda) g^{T}(y - z)$$

$$= \lambda f(x) + (l \cdot \lambda) f(z) = f(z) + g^{T}(\lambda x - \lambda z + (l \cdot \lambda) y - (1 - \lambda) z).$$

$$\bigcirc$$

$$\begin{array}{l} (Onve \times => (\mathbf{x}) \\ \mbox{For } x, y \in dom(f), \quad O < \times < (, \\ ((-\lambda) + [\mathbf{x}] + \lambda + (y) \geq f((-\lambda) \times + \lambda y) = f(\mathbf{x} + \lambda (y - \mathbf{x})) \\ => f(x) \geq f(x) + \frac{f(\mathbf{x} + \lambda (y - \mathbf{x})) - f(\mathbf{x})}{\lambda} \end{array}$$

taking lim as 
$$\lambda \rightarrow 0$$
 gives  
 $f(x) + Df(H^T(\gamma - x))$ 

Thin2 If f is twice differentiable,  
f is convex iff 
$$\nabla^2 f(X)$$
 is psd  
where  $\nabla^2 f(X)$  is matrix with entries  $\overline{D(X)} \overline{D(X)}$  f  
Compating  $\|A\|_2^2 = \max_{\|X\|=1} \|AX\|_2^2 = \max_{\|X\|=1} x^T (A^T A) X$   
 $\|X\|=1$   $\|X\|=1$   $\|X\|=1$   
Does not seem a convex program, because  
*i.* are maximizing  
*ii.*  $\|A\|=1$  is not a convex set.  
Restricting to  $\|\|X\| \le 1$  solves  $\overline{ii}$  but not  $\overline{i}$ .  
Solution: write as  
min t s.t.  $tI - A \ge 0$   
 $(M \ge 0$  iff M is positive semidefinite).  
Is a convex cone.  
(an check if  $M \ge 0$  by trying to compute  
a cholesky factorization : L st.  $LL^T=M$ .

Sn = set of symmetric non positive semidefinite.

If A is symmetric, but 
$$A \notin S_{t,1}^{h}$$
 in space of symmetric  
 $\exists v st. v^{T}Av < O$   
A hyperplane separating A from  $S_{t}^{n}$  is given by  
 $\{ \text{symmetric } X : v^{T}Xv = O \}$   
 $v^{T}Xv = O$  for  $X \in S_{t}^{n}$ .  $v^{T}Av < O$ .  
And is a hyperplane because

$$v^T X v = \sum_{\substack{(z,j) \in N}} X(z,j) v(z) v(z) is linear in X.$$