

Dan's Favorite LP:

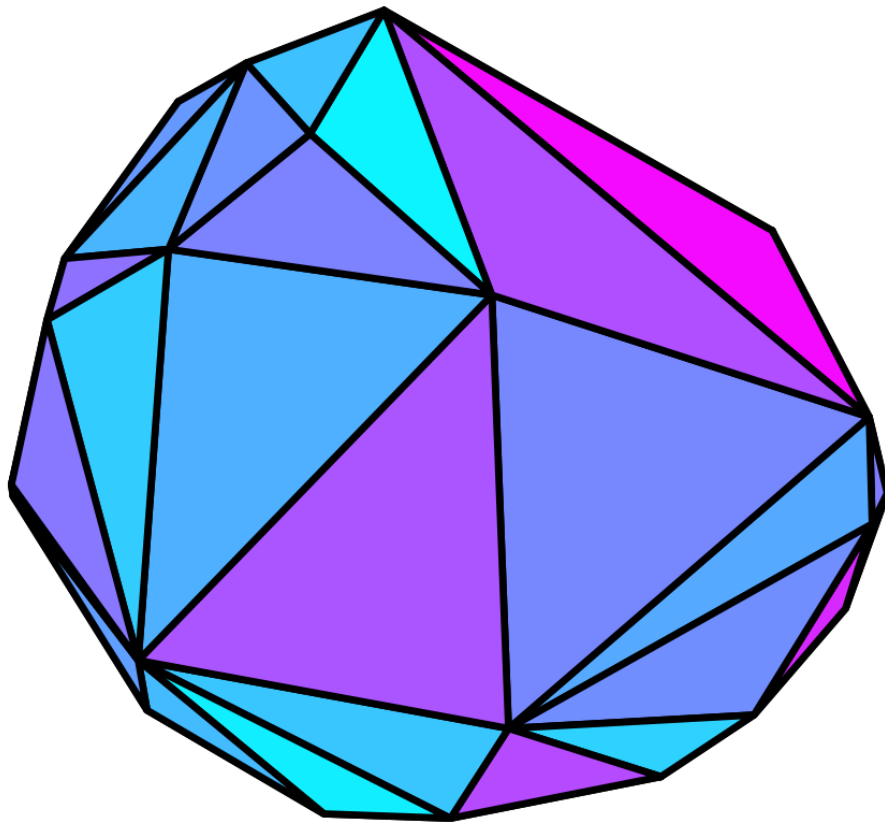
Given  $a_1, \dots, a_m \in \mathbb{R}^d$  st.  $0 \in \text{CH}(a_1, \dots, a_m)$

and  $c \in \mathbb{R}^d$

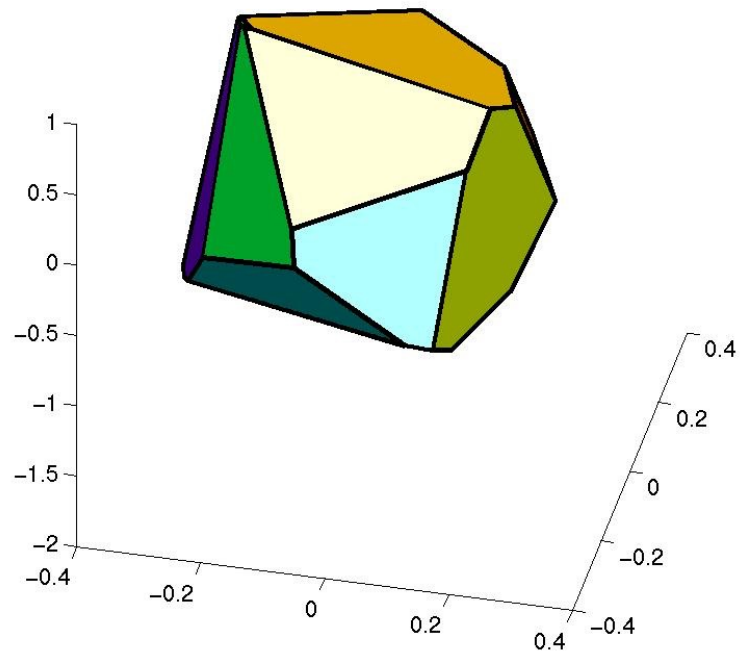
max  $\alpha$  st.  $\alpha c \in \text{CH}(a_1, \dots, a_m)$

Here is a picture of the convex hull of  
random unit vectors in  $\mathbb{R}^3$

If I choose Gaussian random vectors, some  
are probably inside the convex hull of the others



Random hyper-planes give polytopes like this:



To avoid strange behavior,

assume  $c, a_1, \dots, a_m$  are in general position  
which we define to mean:

- i. every subset of  $d$  is linearly independent, and
- ii. every subset of  $d+1$  is affinely independent.  
(so no vector is in the affine span of  $d$  others)

This is true if add a slight random perturbation.

$a_1, \dots, a_d$  are linearly dependent iff  $\sigma_d(a_1, \dots, a_d) = 0$ ,  
and this happens with probability 0 under  
Gaussian noise.

And, the probability one lies on the hyperplane  
defined by the affine span of  $d$  others = 0.

The definition of "General Position" depends  
on the context - so don't get too attached  
to one definition.

The consequence we want is that whenever the  
ray in direction  $c$  intersects a simplex  $\text{CH}(a_i : i \in S)$   
it intersects the relative interior of that simplex



Carathéodory's Theorem If  $\gamma \in \text{CH}(a_1, \dots, a_m)$

then  $\exists |S| \leq d+1$  st.  $\gamma \in \text{CH}(a_i : i \in S)$

proof. Recall  $\gamma \in \text{CH}(a_1, \dots, a_m) \Leftrightarrow \exists t \in \mathbb{R}^m, \mathbb{1}^T t = 1, t \geq 0$   
st.  $\gamma = \sum t(i) a_i$ .

Assume, wolog,  $t(i) > 0$  for all  $i$  and  $m > d+1$

Will show can modify  $t$  to make an entry 0.

As  $m > d+1$ , the vectors  $\begin{pmatrix} a_1 \\ \vdots \\ a_m \end{pmatrix}$  are

linearly dependent. So,  $\exists r \in \mathbb{R}^m, r \neq \bar{0},$  st.

$$\sum r(i) a_i = \bar{0} \quad \text{and} \quad \sum r(i) = 0.$$

So, for every  $\mu,$   $\sum t(i) + \mu r(i) = 1$

$$\text{and} \quad \sum (t(i) + \mu r(i)) a_i = \gamma.$$

Will show can pick  $\mu$  st.  $t + \mu r \geq 0$  (1)

and  $(t + \mu r)(i) = 0$  for some  $i$ . (2)

$\Rightarrow \gamma$  is in CH of  $m-1$  of the vectors.

(1) holds when  $\mu = 0$ .

As  $r \neq \bar{0}$ , is some very positive or very negative  $\mu$   
for which (1) does not hold.

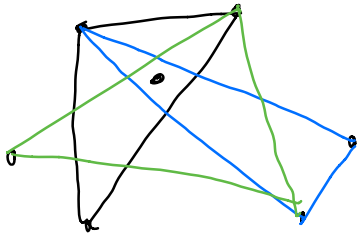
Choose  $\mu$  to be (largest (or smallest)  $\mu$  for  
which (1) holds.

For this  $\mu,$   $t + \mu r$  has a zero entry  $\Rightarrow$  (2)

Note: this proof is algorithmic.

Given  $t$  we can efficiently find  $r$  and  $\mu$ .

The set  $S$  need not be unique:



and are many more

The Dual LP:

Geometric view: if  $\mathcal{H} = \{b : y^T b = 1\}$  is a hyperplane s.t. all of  $a_1, \dots, a_m$  is on one side, and  $\alpha c \in \mathcal{H}$ , then value of LP  $\leq \alpha$ .

$$\alpha c \in \mathcal{H} \Leftrightarrow y^T(\alpha c) = 1 \Rightarrow \alpha = 1/y^T c$$

So, LP is  $\max y^T c$  s.t.  $y^T a_i \leq 1, \forall i$

We just established

Weak Duality: value of dual upper bounds value of primal.

Algebraically: if  $\alpha c = \sum_i t(i) a_i, \sum_i t(i) = 1,$

$$y^T(\alpha c) = \sum t(i) y^T a_i \leq \sum t(i) = 1$$

Strong Duality: Optima of both are the same.

That is, if  $\alpha_* = \max \alpha$  s.t.  $\alpha c \in \text{CH}(a_1, \dots, a_m)$

and  $\gamma_*$  maximizes  $\gamma^T c$  s.t.  $\gamma^T a_i \leq 1$

Then  $\gamma_*^T(\alpha_* c) = 1$ .

proof Let the supporting hyperplane of  $\text{CH}(a_1, \dots, a_m)$  at  $\alpha_* c$  be  $\mathcal{H} = \{b : \gamma^T b = 1\}$ .

So,  $\gamma^T a_i \leq 1$  for all  $i$  and  $\gamma^T(\alpha_* c) = 1$

The Simplex Method, or how to find  $\max \alpha$  s.t.

$\alpha c \in \text{CH}(a_1, \dots, a_m)$ . Will involve solving many systems of equations. Will not have useful run-time guarantees.

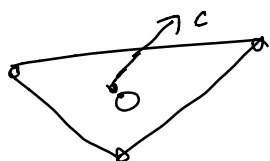
Start: Assume that  $0 \in \text{interior}(\text{CH}(a_1, \dots, a_m))$ ,

AND that we know  $t$  s.t.  $\bar{0} = \sum_i t(i) a_i$

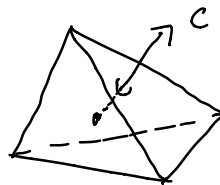
(We will later discuss how we find this  $t$ )

Find such a  $t$  that has  $d+1$  nonzero coordinates.

So,  $\bar{0} \in \Delta$  a simplex.



or



wolog, let  $a_1, \dots, a_{d+1}$  be the corners

First step: find the point on the boundary hit by the ray in direction  $c$ . That is,

$$\max \alpha \text{ s.t. } \alpha c \in \text{CH}(a_1, \dots, a_{d+1}).$$

There are only  $d+1$  faces to check

(each obtained by dropping a vertex).

If that face is  $\text{CH}(a_1, \dots, a_d)$ ,

then there is an  $\alpha$  s.t.  $\alpha c \in \text{CH}(a_1, \dots, a_d)$ .

To find it, solve the equation

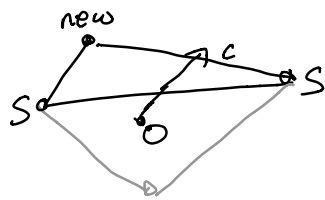
$$c = \sum_{i=1}^d x(i) a_i$$

It is the right face if all  $x(i) \geq 0$

Let  $\alpha = (\sum x(i))^{-1}$  and  $t(i) = \alpha x(i)$

$$\text{So, } \alpha c = \sum_{i=1}^d t(i) a_i.$$

General step: begin knowing  $\alpha c = \sum_{i \in S} t(i) a_i$  for some  $|S|=d$ .



- a  $(d-1)$ -dim simplex.

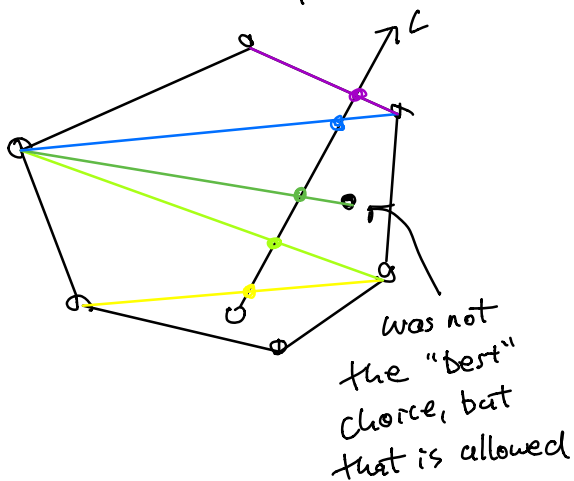
Find an  $a_j$  s.t.  $\hat{\alpha} c \in \text{CH}(a_j \cup (a_i)_{i \in S})$ , for  $\hat{\alpha} > \alpha$   
that is a simplex that contains the ray on the other side of the  $(d-1)$ -dim simplex.

We then proceed as before to find the point where the ray leaves this new simplex, and again express it as a positive sum of  $d$  of the points  $a_i$ .

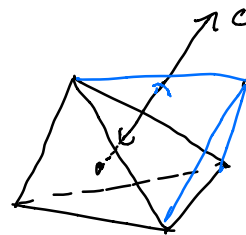
Note: our assumptions about "General Position" imply that  $\alpha c$  will always be in the interior of the  $(d-1)$  simplex, and thus  $(d+\varepsilon)c$  will be in the interior of the  $d$ -simplex, for small  $\varepsilon > 0$ .

We use the fact that a ray that goes through the interior of a simplex intersects its boundary at two points: where it enters and where it leaves.

A 2D example:



3D





When does it stop?

When there is no simplex on the other side.

Let the current point be  $\alpha C$  and  $S \subseteq \{1..m\}$ ,  $|S|=d$   
index the simplex containing it.

$\alpha C \in$  relative interior  $\text{CH}(a_i : i \in S)$

Claim: no simplex on other side  $\Rightarrow \alpha C$  on boundary of tope.

Let  $H =$  affine span  $(a_i : i \in S)$   
 $= \{z : \gamma^T z = 1\}$  for some  $\gamma$ .

We have  $\gamma^T a_i = 1$   $i \in S$

and  $\gamma^T a_i \leq 1$  for  $i \notin S$ .

So,  $H$  is a supporting hyperplane of  $\text{CH}(a_1, \dots, a_m)$  at  $\alpha C$ ,

$\Rightarrow \alpha$  is the optimum

The Dual simplex method:

Find  $\gamma$  maximizing  $\gamma^T C$  s.t.  $\gamma^T a_i \leq 1 \ \forall i$ .

This is the dual LP:  $\max \gamma^T C$   
 $\gamma^T a_i \leq 1, \ \forall i$

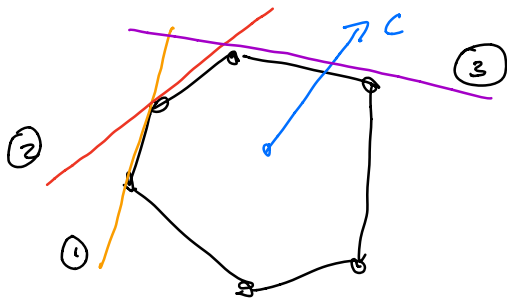
Imagine that know some supporting

hyperplane:  $\gamma$  s.t.  $\gamma^T a_i = 1 \ i \in S, \ |S|=d$

and  $\gamma^T a_i \leq 1 \ \forall i$

Want to rotate it around the polytope to find a supporting hyperplane that increases  $\gamma^*C$ .

In 2D, looks like this



To see is always possible, if  $\gamma$  is not optimal, look at intersection of  $\alpha C$  in that plane, and the simplex:

