

Sparse recovery 2: Recovering sparse signals with noise.

Problem: Given d -by- n A , $d \ll n$ $d \left(\begin{matrix} n \\ A \end{matrix} \right)$
and a vector $b \in \mathbb{R}^d$, try to find a
sparse x s.t. $Ax \approx b$.

Formally:

$$\text{Given } \underline{\varepsilon} > 0, \quad \min_{\substack{x_0 \\ \|x\|_0}} \|x\|_0 \quad \text{s.t.} \quad \underline{\|Ax - b\|_2} \leq \varepsilon \quad (P_0)$$

Is NP-hard, but can approximately solve it for nice A .

$$\text{Let } x_* \text{ be the } \begin{matrix} \text{convex} \\ \text{solution to} \end{matrix} \quad \min \underline{\|x\|_1} \quad \text{s.t.} \quad \|Ax - b\|_2 \leq \varepsilon. \quad \begin{matrix} \text{convex} \\ (P_1) \end{matrix}$$

We will show x_* is close to the solution to P_0 .

$$\begin{aligned} d &= \text{const} \cdot n \\ s &= \text{const} \cdot d \end{aligned}$$

$$\text{as } \frac{s}{d} \rightarrow 0, \quad \frac{d}{s} \text{ is fixed} \\ \delta \rightarrow 0$$

Nice: A satisfies the (s, δ) -Restricted Isometry Property

if for all x s.t. $\|x\|_0 \leq s$

$$\underline{(1-\delta)\|x\|_2} \leq \|Ax\|_2 \leq \underline{(1+\delta)\|x\|_2}$$

\leftarrow usually are squares

Def $\delta_s(A) =$ smallest δ s.t. A satisfies (s, δ) -RIP

Theorem If $\delta_{3s} + \frac{1}{\sqrt{2}} \delta_{2s} \leq 1 - \sqrt{2}$, $b = Ax_0 + e$

with $\|x_0\|_0 \leq s$ and $\|e\|_2 \leq \varepsilon$,

then x_* , the solution to P_1 , satisfies

$$\|x_0 - x_*\|_2 \leq \frac{2\sqrt{2} \varepsilon}{1 - \sqrt{2} - \delta_{3s} - \frac{1}{\sqrt{2}} \delta_{2s}} \text{ const}$$

Notation: let $S = \text{supp}(x_0)$, $s = |S|$, $T = \{1, \dots, n\} - S$.

$$h = x_* - x_0 \quad \|h\|_2 < ? \quad x_0(i) = 0 \quad i \in T$$

Claim 1 $\|Ah\|_2 \leq 2\varepsilon$ USE RIP $\|h\|_2 \leq ? \|Ah\|_2$

proof $Ah = Ax_* - Ax_0 = (Ax_* - b) - (Ax_0 - b)$

$$\text{so } \|Ah\|_2 \leq \|Ax_* - b\|_2 + \|Ax_0 - b\|_2 \leq 2\varepsilon$$

$< \varepsilon \qquad \qquad \qquad < \varepsilon$

Let $h(S)$ be the restriction of h to coordinates in S .
(or just zero out coordinates not in S)

Claim 2 $\|h(S)\|_1 \geq \|h(T)\|_1$ $T = \{1, \dots, n\} - S$

proof $\|x_0\|_1 \geq \|x_*\|_1$ because x_* min $\|x\|_1$ s.t.

$$= \|x_0 + h\|_1 = \|(x_0 + h)(S)\|_1 + \|(x_0 + h)(T)\|_1$$

$\swarrow \Delta - \text{neg} \qquad \qquad \qquad \downarrow x_0(T) = 0$

$$\geq \|x_0(S)\|_1 - \|h(S)\|_1 + \|h(T)\|_1$$

$$= \|x_0\|_1 - \|h(S)\|_1 + \|h(T)\|_1$$

$$\Rightarrow \|h(S)\|_1 \geq \|h(T)\|_1$$

Would also like to say $\|h(S)\|_2 \geq \text{const} \cdot \|h(T)\|_2$,

but might not be true.

Ex. $h(S) = \underbrace{\frac{1}{s}, \dots, \frac{1}{s}}_{s \text{ coords}}$ $\|h(S)\|_1 = 1$, $\|h(S)\|_2 = \frac{1}{\sqrt{s}}$
 $h(T) = 1, 0, 0, \dots, 0$ $\|h(T)\|_1 = 1$, $\|h(T)\|_2 = 1$

Will show something like this after excluding $2s$ coordinates from T .

Def. $T_i =$ the $2s$ coords in T on which h is largest

$$R = T - T_i = \{1, \dots, n\} - S - T_i$$

Claim 3 $\|h(R)\|_2^2 \leq \frac{1}{2} \|h(S)\|_2^2$ (only uses claim 2)

proof

$$\|v\|_2^2 \leq \|v\|_\infty \|v\|_1$$

$$\Leftrightarrow \sum_i v(i)^2 \leq \sum_i |v(i)| \cdot \|v\|_\infty = \|v\|_1 \cdot \|v\|_\infty$$

$$\|h(R)\|_\infty \leq \min h(T_i) \leq \frac{1}{2s} \|h(T_i)\|_1 \quad (\text{min} \leq \text{average})$$
$$\leq \frac{1}{2s} \|h(T)\|_1, \quad T_i \subset T$$

$$\|h(R)\|_1 \leq \|h(T)\|_1, \quad R \subset T$$

$$\|h(R)\|_2^2 \leq \frac{1}{2s} \|h(T)\|_1^2 \leq \frac{1}{2s} \|h(S)\|_1^2$$
$$\|v\|_1 \leq \sqrt{\|v\|_0} \|v\|_2, \quad \forall v$$
$$\leq \frac{s}{2s} \|h(S)\|_2^2$$
$$= \frac{1}{2} \|h(S)\|_2^2$$

We also want to show $\|A h(R)\|_2 \leq \text{something}$.

Will need RIP, but only applies to sparse vectors.

$T_1 = 2s$ coords of T on which h is largest

Let $T_2 =$ largest $2s$ coords of h in $T - T_1$

$T_3 =$ " " " " " " $T - T_1 - T_2$

etc.

So, $T = \underline{T_1 \cup T_2 \cup \dots \cup T_k}$, some k , and T_k can be smaller.

$R = \underline{T_2 \cup \dots \cup T_k}$ $\{1, \dots, n\} = S \cup T_1 \cup R$

Claim 4 $\sum_{i=2}^k \underline{\|h(T_i)\|_2} \leq \underline{\frac{1}{\sqrt{2}} \|h(S)\|_2}$ (w/ $\|h\|_0 \leq ?$)

proof For $i \geq 2$, $\|h(T_i)\|_2^2 \leq \|h(T_i)\|_\infty \|h(T_i)\|_1$
 $\leq \|h(T_i)\|_\infty \|h(T_{i-1})\|_1$ h decreases as $i \uparrow$
 $\leq \frac{1}{2s} \|h(T_{i-1})\|_1^2$ $\|h(T_i)\|_\infty \leq \max_{h(T_{i-1})}$

$$\Rightarrow \sum_{i=2}^k \|h(T_i)\|_2 \leq \frac{1}{\sqrt{2s}} \sum_{i=1}^{k-1} \|h(T_i)\|_1 \quad \frac{1}{2s} \|h(T_{i-1})\|_1$$

$$= \frac{1}{\sqrt{2s}} \|h(T_1 \cup \dots \cup T_{k-1})\|_1 \leq \frac{1}{\sqrt{2s}} \|h(T)\|_1$$

$$\leq \frac{1}{\sqrt{2s}} \|h(S)\|_1 \quad (\text{Claim 2})$$

$$\leq \frac{1}{\sqrt{2}} \|h(S)\|_2, \text{ as } \|h(S)\|_0 \leq S$$

Claim 5 $\|A h(R)\|_2 \leq \frac{1}{\sqrt{2}} (1 + \delta_{2s}) \|h(S)\|_2$

$$\|A h(R)\|_2 \stackrel{\Delta}{\leq} \sum_{i=2}^k \|A h(T_i)\|_2 \leq (1 + \delta_{2s}) \sum_{i=2}^k \|h(T_i)\|_2$$

$$\leq (1 + \delta_{2s}) \frac{1}{\sqrt{2}} \|h(S)\|_2$$

proof of theorem $\{1, \dots, n\} = S \cup T_1 \cup R$

$$\|h\|_2 = \|x_0 - x_*\|_2 \leq \frac{2\sqrt{2}\varepsilon}{1 - \sqrt{2} - \sigma_{3s} - \frac{1}{\sqrt{2}}\sigma_{2s}}$$

$$\|Ah\|_2 \leq 2\varepsilon$$

$$\|Ah(R)\|_2 \leq \frac{1}{\sqrt{2}}(1 + \sigma_{2s})\|h(S)\|_2 \leq \frac{1}{\sqrt{2}}(1 + \sigma_{2s})\|h(S \cup T_1)\|_2$$

$$\Delta\text{-ineq} \quad 2\varepsilon \geq \|Ah\|_2 \geq \|Ah(S \cup T_1)\|_2 - \|Ah(R)\|_2$$

$$|S \cup T_1| = 3s$$

$$\geq (1 - \sigma_{3s})\|h(S \cup T_1)\|_2 - \frac{1}{\sqrt{2}}(1 + \sigma_{2s})\|h(S \cup T_1)\|_2$$

$$= \left[1 - \sqrt{2} - \sigma_{3s} - \frac{1}{\sqrt{2}}\sigma_{2s}\right]\|h(S \cup T_1)\|_2$$

$$\begin{aligned}\|h(S \cup T_1)\|_2^2 &= \|h\|_2^2 - \|h(R)\|_2^2 \\ &\geq \|h\|_2^2 - \frac{1}{2}\|h(S)\|_2^2 \geq \frac{1}{2}\|h\|_2^2\end{aligned}$$

$$2\varepsilon \geq \left[1 - \sqrt{2} - \sigma_{3s} - \frac{1}{\sqrt{2}}\sigma_{2s}\right] \frac{1}{\sqrt{2}}\|h\|_2$$